

Weighted finite-state automata

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Finite-state automata

Motivating example: a spell checker

The language model: a weighted automaton

The spelling model: a transducer

The error model: a weighted transducer

Parameter estimation

Rational kernels

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Finite-state acceptor

A **finite-state acceptor** accepts sequences of symbols that correspond to paths from a start state to a final state.



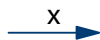
State (label has no meaning)



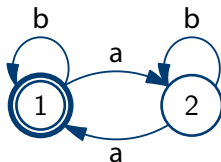
Start state



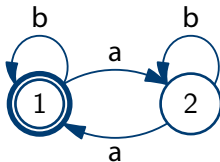
Final state



Transition that accepts symbol x



This is **not** a graphical model. It is blue.



We will see

- **weighted** automata;
- **n -tape** automata.
 - $n = 2$: **transducers**.

Motivating example: a spell checker

This **is** a graphical model. It is orange and has factors.



with

- W the words;
- C the correct characters;
- K the key presses;

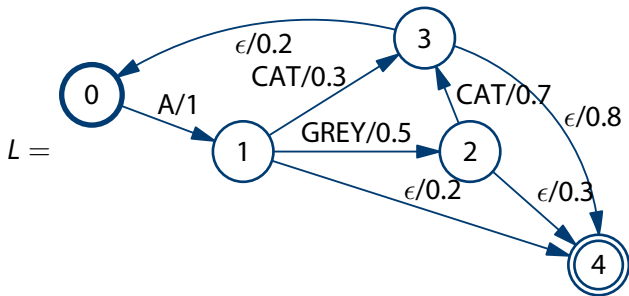
and

- L the language model;
- S the spelling model;
- T the error (typo) model.

The language model L : a weighted automaton

n -gram model:

$$P(W) \simeq P(w_1)P(w_2|w_1)P(w_3|w_2)\cdots P(w_n|w_{n-1})$$
$$\triangleq [[L]](W)$$

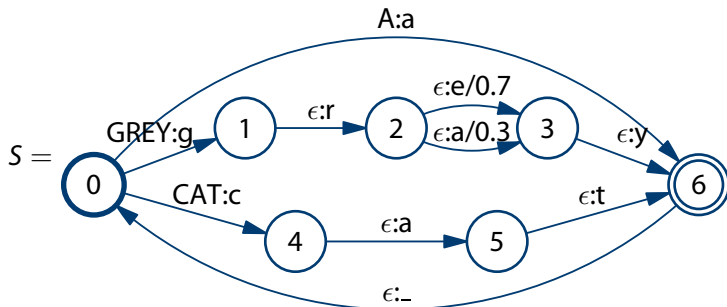


- /0.3 indicates a weight of 0.3.
- ϵ means no symbol.

The spelling model: a transducer

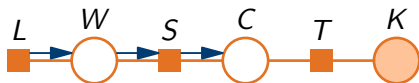


$$P(C|W) = [[S]](W, C)$$



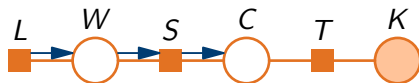
- a:b indicates transduction from a to b.

Message passing

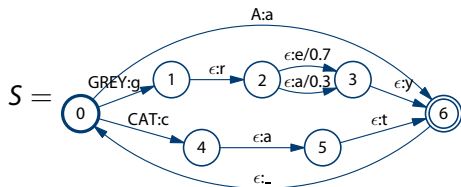
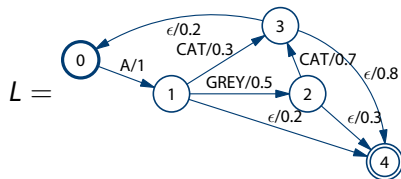


What are the messages?

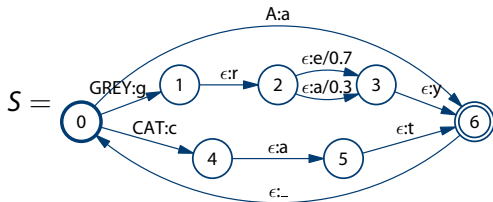
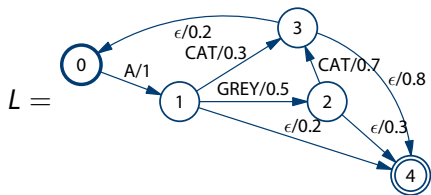
Message passing



What are the messages?



Composition

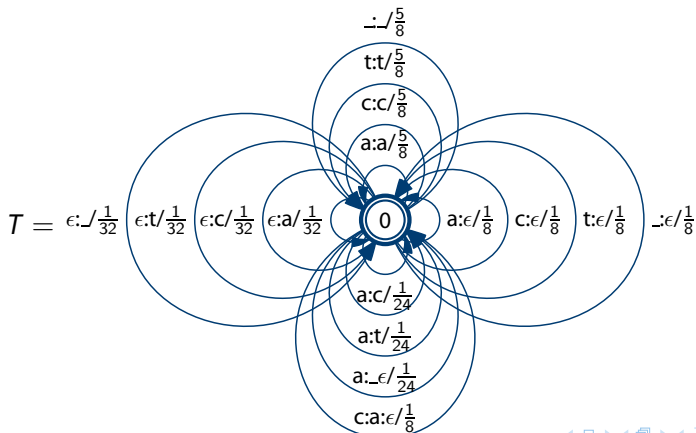


- Composition: $L \circ S$;
- projection.

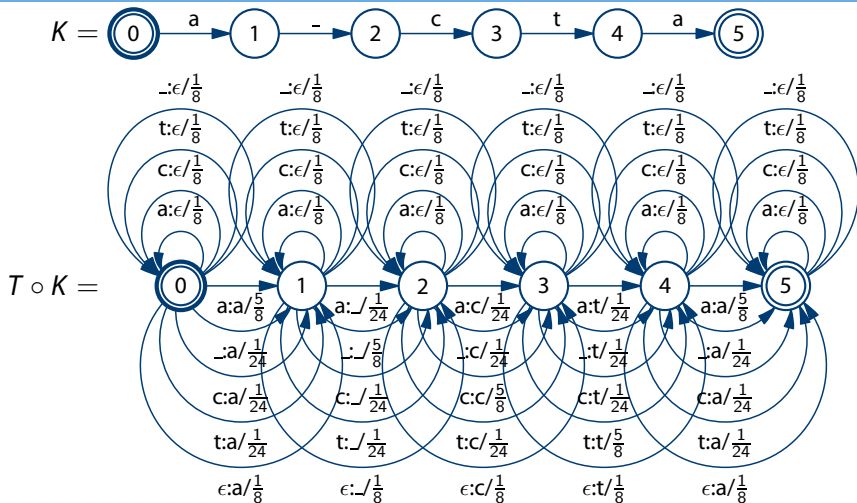
The error model: a weighted transducer



$$P(K|C) \propto [[T]](C, K)$$



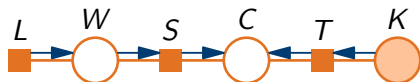
Composition



What does it mean to have multiple paths with input “a _ c a t”?



Different semirings for weights



- $L \circ S \circ T \circ K$ scales linearly with the input.
- $\max_C [[L \circ S \circ T \circ K]] (C)$ means summing over all paths.
- Instead, take the best path.
- Formally, use a different semiring for the weights:
 - $(\mathbb{K}, \oplus, \otimes, \bar{0}, \bar{1}) = (\mathbb{R}, +, \times, 0, 1)$;
 - $(\mathbb{K}, \oplus, \otimes, \bar{0}, \bar{1}) = (\mathbb{R}, \max, \times, 0, 1)$;
 - and their logarithms (e.g. “tropical semiring”).
- Always use a “shortest-distance” algorithm.
- Roughly equivalent to forms of message-passing:
 - sum-product algorithm;
 - max-product algorithm;
 - max-sum algorithm.

- L and S are known beforehand.



- So they can be composed off-line:



- $L \circ S$ can be optimised.
 - Determinisation/minimisation need to know what \oplus means.

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Parameter estimation

Rational kernels

- If our parameters are the weights of an automaton, how to estimate them?
- “Expectation semiring” (Eisner 2002).

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A kernel K is **rational** (Cortes 2004) if it can be defined as

$$K(x, y) = \psi(\llbracket T \rrbracket(x, y))$$

where

- T is a transducer over some semiring \mathbb{K} ; and
- $\psi : \mathbb{K} \rightarrow \mathbb{R}$.

If

- T is a sensible transducer over some semiring \mathbb{K} ; and
- $\psi : \mathbb{K} \rightarrow \mathbb{R}$ is a semiring morphism,

then

$$K(x, y) = \psi([[T \circ T^{-1}]](x, y))$$

is positive definite and symmetric.

Weighted finite-state automata:

- a useful formalism for models over symbol sequences.
- For inference.
- For string kernels.